



# Daffodil International University

Faculty of Science & Information Technology

Department of Computer Science & Engineering

Mid Examination, Fall 2025

Course Code: CSE313, Course Title: Compiler Design

Level: 3 Term: 1 Batch: 65

Time: 01:30 Hrs

Marks: 25

Answer ALL Questions

*[The figures in the right margin indicate the full marks and corresponding course outcomes. All portions of each question must be answered sequentially.]*

| 1.    | a)            | <p><math>S \rightarrow \text{WHP ADJ PUNCT} \mid \text{WH INTENS ADJ PUNCT}</math></p> <p>WHP <math>\rightarrow</math> why so</p> <p>WH <math>\rightarrow</math> why</p> <p>INTENS <math>\rightarrow</math> so</p> <p>ADJ <math>\rightarrow</math> serious   angry   calm</p> <p>PUNCT <math>\rightarrow</math> !   ?</p> <p>For the input command: “why so serious!”, briefly explain whether the given CFG is ambiguous or unambiguous.</p> | [5]   | CO1           |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
|-------|---------------|-----------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------|-------|---------------|---------------|----|----------|------|----|------|-------------|----|-------------|------|----|------|------|--|-----|
| 2.    | a)            | <p>MoodMeter ( <i>float joy, float sarcasm, float caffeine</i> ) {</p> <p>float 1st_Score;</p> <p>Score = joy * Level - sarcasm / 2 + caffeine * 0.5     // <b>target statement</b></p> <p>print(“ %f”, 1st_Score);</p> <p>}</p> <p><b>Describe</b> the phases of a compiler for the <b>target statement</b> of the given function.</p>                                                                                                       | [6]   |               | CO1           |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
|       | b)            | <p><b>Find</b> the type of errors in the above code 2(a) and <b>Why</b> is Symbol Table Correction considered a better error recovery method in semantic analysis, and is there any strategy that works better in certain cases? Justify your answer.</p>                                                                                                                                                                                     | [4]   |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
| 3.    |               | <p>Consider an NFA with the following transition table:</p> <table><tr><th>State</th><th><math>\delta(q,a)</math></th><th><math>\delta(q,b)</math></th></tr><tr><td>q0</td><td>{q0, q1}</td><td>{q0}</td></tr><tr><td>q1</td><td>{q2}</td><td><math>\emptyset</math></td></tr><tr><td>q2</td><td><math>\emptyset</math></td><td>{q3}</td></tr><tr><td>q3</td><td>{q3}</td><td>{q3}</td></tr></table>                                          | State | $\delta(q,a)$ | $\delta(q,b)$ | q0 | {q0, q1} | {q0} | q1 | {q2} | $\emptyset$ | q2 | $\emptyset$ | {q3} | q3 | {q3} | {q3} |  | CO1 |
| State | $\delta(q,a)$ | $\delta(q,b)$                                                                                                                                                                                                                                                                                                                                                                                                                                 |       |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
| q0    | {q0, q1}      | {q0}                                                                                                                                                                                                                                                                                                                                                                                                                                          |       |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
| q1    | {q2}          | $\emptyset$                                                                                                                                                                                                                                                                                                                                                                                                                                   |       |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
| q2    | $\emptyset$   | {q3}                                                                                                                                                                                                                                                                                                                                                                                                                                          |       |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
| q3    | {q3}          | {q3}                                                                                                                                                                                                                                                                                                                                                                                                                                          |       |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
|       | a)            | <p>Apply the <b>subset construction method</b> to convert a given NFA into an equivalent DFA. Show the <b>complete subset construction table</b> with all reachable states and draw the <b>state diagram</b> of the resulting DFA. Identify and <b>justify</b> which states in the DFA are the <b>final states</b>.</p>                                                                                                                       | [5]   |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |
|       | b)            | <p>Write a <b>regular expression</b> that describes the language accepted by the resulting DFA. Furthermore, <b>explain</b> the relationship between the number of states in the original NFA and the constructed DFA (i.e., why the DFA has fewer, more, or the same number of states) and identify the <b>three main cases of non-determinism</b> that were eliminated during this conversion process.</p>                                  | [5]   |               |               |    |          |      |    |      |             |    |             |      |    |      |      |  |     |