



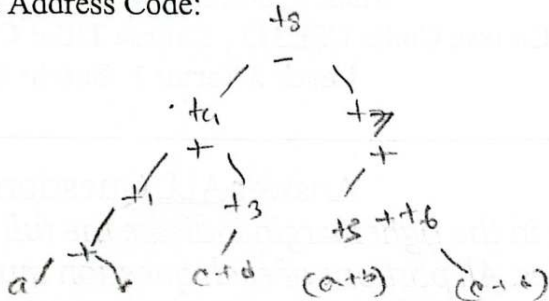
Daffodil International University
Faculty of Science & Information Technology
Department of Computer Science & Engineering
Final Examination, Spring 2026
Course Code: CSE313 , Course Title: Compiler Design
Level: 3 Term: 1 Batch: 66

Time: 02:00 Hrs
Marks: 40

Answer ALL Questions

[The figures in the right margin indicate the full marks and corresponding course outcomes. All portions of each question must be answered sequentially.]

1.	a)	Consider the following grammar to produce LR(0) parser and canonical table and identify any shift-reduce conflicts, if present. $S \rightarrow X + Y Z$ $Y \rightarrow W n \mid n \mid W$ $X \rightarrow p \mid q$ $W \rightarrow r \mid s \mid t$ $Z \rightarrow u \mid v$	[8]	CO2
	b)	Consider the following code segment: <pre>sum = 0; for (i = 0; i < n; i++) { x = a * b; y = i * 2; sum = sum + x + y; }</pre> Apply appropriate code optimization techniques to optimize the given code and rewrite the optimized code using suitable methods and justify the applied techniques.	[2]	
2.	a)	Given the grammar below, apply predictive parsing techniques to determine the FIRST and FOLLOW sets of each nonterminal. Then, use these sets to build the LL(1) parsing table. $P \rightarrow S Q$ $S \rightarrow S a \mid b T \mid \epsilon$ $T \rightarrow c T \mid d$ $Q \rightarrow f Q \mid g$	[7]	CO2
	b)	Given the grammar below, eliminate left factoring and produce an equivalent left-factored grammar: $S \rightarrow \text{stmt} \mid \text{stmt block} ; \text{stmt} \mid \text{stmt block} , \text{stmt}$ $\text{stmt} \rightarrow \text{while expr do stmt} \mid \text{while expr do stmt block}$	[3]	

3.	a)	Consider the following arithmetic expression: $(a + b) * (c - d) + (a + b) / (-e) + (f - d)$ Apply intermediate code generation techniques to evaluate the expression. i. Apply your knowledge to write the Three Address Code (TAC) for the expression. ii. Generate the corresponding Quadruples representation. iii. Analyze the structure and produce the Indirect Triples representation where you have 25, 27, 29, 31, 33, 35 and 41 available memory in STACK.	[6]	CO3																																										
	b)	Consider the following Three Address Code: 1. $t1 = a + b$ 2. $t3 = c + d$ 3. $t4 = t1 + t3$ 4. $t5 = a + b$ 5. $t6 = c + d$ 6. $t7 = t5 + t6$ 7. $t8 = t4 - t7$ 	[4]																																											
4.	a)	Consider the following Three Address Code: <table border="1"><tr><td>1. $i = 0$</td><td>15. $t6 = 4 * i$</td><td>29. $t12 = 4 * j$</td></tr><tr><td>2. $j = n$</td><td>16. $y = a[t6]$</td><td>30. $p = a[t12]$</td></tr><tr><td>3. $k = m$</td><td>17. $t7 = 4 * j$</td><td>31. If $p == x$ goto (34)</td></tr><tr><td>4. $t1 = 4 * i$</td><td>18. $t8 = a[t7]$</td><td>32. $q = p + x$</td></tr><tr><td>5. $x = a[t1]$</td><td>19. $a[t6] = t8$</td><td>33. goto (36)</td></tr><tr><td>6. $i = i + 1$</td><td>20. $a[t7] = y$</td><td>34. $q = p - x$</td></tr><tr><td>7. $t2 = 4 * i$</td><td>21. $k = k - 1$</td><td>35. $r = q * 2$</td></tr><tr><td>8. $t3 = a[t2]$</td><td>22. If $k > 0$ goto (6)</td><td>36. $s = r + 1$</td></tr><tr><td>9. If $t3 < x$ goto (6)</td><td>23. $t9 = 4 * i$</td><td>37. If $s < 100$ goto (21)</td></tr><tr><td>10. $j = j - 1$</td><td>24. $z = a[t9]$</td><td>38. $t13 = 4 * k$</td></tr><tr><td>11. $t4 = 4 * j$</td><td>25. $t10 = 4 * n$</td><td>39. $a[t13] = s$</td></tr><tr><td>12. $t5 = a[t4]$</td><td>26. $t11 = a[t10]$</td><td>40. end</td></tr><tr><td>13. If $t5 > x$ goto (10)</td><td>27. $a[t9] = t11$</td><td></td></tr><tr><td>14. If $i >= j$ goto (25)</td><td>28. $a[t10] = z$</td><td></td></tr></table> Answer the following question: i. Which lines in the given three-address code are identified as leaders? ii. Is / Are there any instruction(s) designated as leaders using both Rule 2 and Rule 3? iii. What is the total number of basic blocks identified in the code? iv. Draw the flow graph for the above-mentioned instructions.	1. $i = 0$	15. $t6 = 4 * i$	29. $t12 = 4 * j$	2. $j = n$	16. $y = a[t6]$	30. $p = a[t12]$	3. $k = m$	17. $t7 = 4 * j$	31. If $p == x$ goto (34)	4. $t1 = 4 * i$	18. $t8 = a[t7]$	32. $q = p + x$	5. $x = a[t1]$	19. $a[t6] = t8$	33. goto (36)	6. $i = i + 1$	20. $a[t7] = y$	34. $q = p - x$	7. $t2 = 4 * i$	21. $k = k - 1$	35. $r = q * 2$	8. $t3 = a[t2]$	22. If $k > 0$ goto (6)	36. $s = r + 1$	9. If $t3 < x$ goto (6)	23. $t9 = 4 * i$	37. If $s < 100$ goto (21)	10. $j = j - 1$	24. $z = a[t9]$	38. $t13 = 4 * k$	11. $t4 = 4 * j$	25. $t10 = 4 * n$	39. $a[t13] = s$	12. $t5 = a[t4]$	26. $t11 = a[t10]$	40. end	13. If $t5 > x$ goto (10)	27. $a[t9] = t11$		14. If $i >= j$ goto (25)	28. $a[t10] = z$		[7]	CO3
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	b)	Explain Dead Code Elimination and Loop Invariants. Provide examples to show how each optimization works.	[3]																																											